GWDG – Kurs Parallel Programming with MPI

MPI Collective Operations

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Learning Objectives

- Why Add Collective Operations to Send-Receive Message Passing
- MPI Functions for Collective Operations:

```
MPI_BCAST, MPI_SCATTER, MPI_GATHER,
MPI REDUVE, MPI BARRIER
```

Outline

- Why Collective Operations?
- Data-Patterns of Collectives
- Syntax of Communcation Collectives:

```
MPI_BARRIER
MPI_BCAST
MPI_SCATTER
MPI_GATHER
and variants
```

Reduction Collectives

```
MPI_REDUCE
MPI_REDUCE_SCATTER
MPI_SCAN, MPI_EXSCAN
```

Motivation for Collective Operations

- Process 0 provides an input value: copy this value to all processes
- All processes compute a number: collect these numbers into an array in process 0
- All processes compute a number: determine the maximum of these numbers

MPI provides single function calls to perform the necessary multiple point-to-point message passing communication for collective communication patterns

Motivation for Collective Operations

The implementation can provide different realizations for collective operations, optimized with respect to the underlying hardware and to the data layout in the application.

For Intel MPI implementation, see

https://software.intel.com/en-us/mpi-developer-reference-linux-i-mpi-adjust-family-environment-variables

For OpenMPI, see

Choosing a Specific Collective Algorithm Implementation in OpenMPI

and

https://www.open-mpi.org/faq/?category=tuning#mca-params

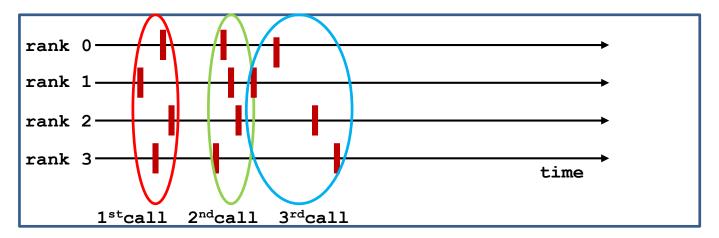
Motivation for Collective Operations

Example: Process 0 copies data to all processes

```
Point-to-point message passing:
 if (me == 0) {
                                destination
   for ( ip=1;ip<np;ip++ ) {
     MPI Send(&val, 1, MPI INT, ip, 0, com );}
                            source
  else
     MPI_Recv(&val, 1, MPI INT, 0, 0, com,
                MPI STATUS IGNORE );
                                    source
                                            no taq
Collective:
MPI Bcast(&val, 1, MPI Type int, 0, com );
```

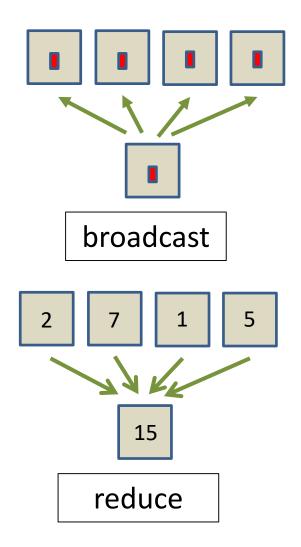
Characeristics of Collective Operations

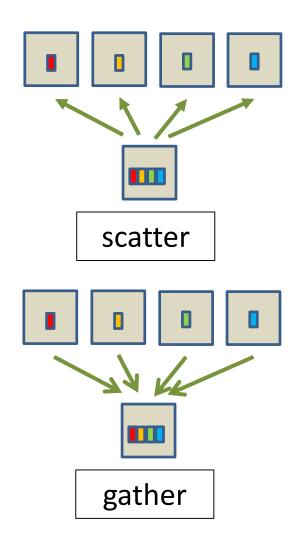
- All processes of a communicator must participate, i.e. must call the collective routine.
- On a given communicator, the n-th collective call must match on all processes of the communicator. Therefore, no tags needed for collective operations.
- If one or more processes of a communicator do not participate in a given collective operation, the program will hang.



- In MPI-1.0 MPI-2.2, all collective operations are blocking.
- Non-blocking versions since MPI-3.0.
- buffers on all processes must have exactly the same size.

Basic Types of Collective Operations





Classification of Collective Operations

MPI BARRIER:

MPI_BCAST:

MPI_GATHER:

MPI SCATTER:

MPI_ALLGATHER:

MPI ALLTOALL:

MPI REDUCE:

MPI_ALLREDUCE:

MPI_REDUCE_SCATTER:

MPI_SCAN, MPI_EXSCAN:

Synchronisation

Send from one process to all processes

gather data from all processes on one process

scatter data from one process to all processes

gather data from all processes, broadcast them

to all processes

exchange data between all processes

reduction over all processes, result goes to

one process

reduction over all processes, result is

broadcasted to all processes

reduction over all processes, result is

scattered to all processes

process i receives result from reduction over

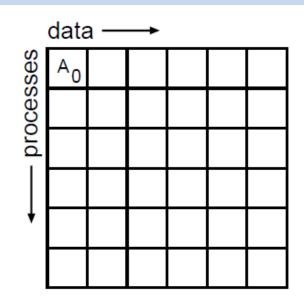
processes with j<= i, j<i

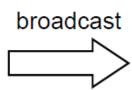
one \rightarrow all

all \rightarrow one

all \rightarrow all

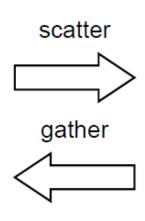
Data Flow in Collective Communication





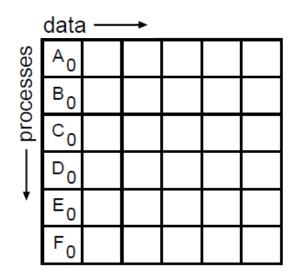
A ₀			
A ₀			

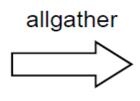
A ₀	A ₁	A ₂	A ₃	A ₄	A ₅



A ₀			
A ₁			
A ₂			
A ₃			
A ₄			
A ₅			

Kollektive Operationen: Datenfluss





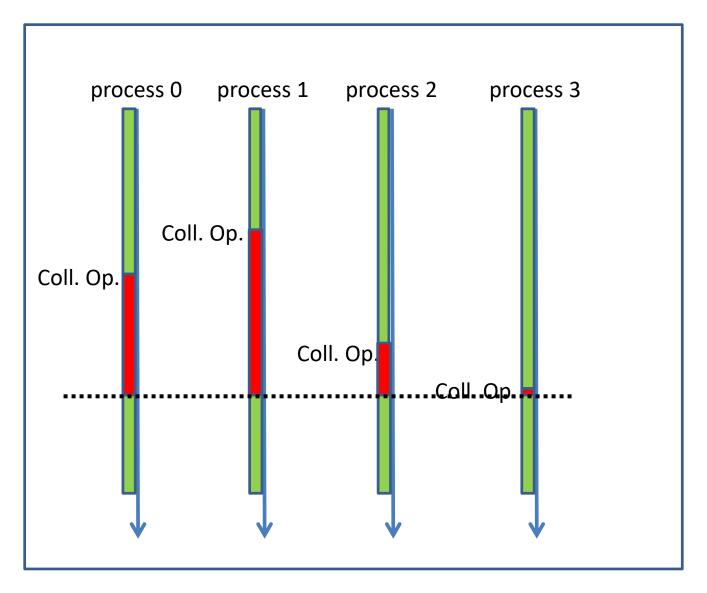
A ₀	В ₀	c ₀	D ₀	E ₀	F ₀
			D ₀	E ₀	F ₀
A ₀	В ₀	c ₀	D ₀	E ₀	F ₀
A ₀	В ₀	C ₀	D ₀	E ₀	F ₀
A ₀	В ₀	c ₀	D ₀	E ₀	F ₀
A ₀	В ₀	c ₀	D ₀	E ₀	F ₀

A ₀	A ₁	A ₂	A ₃	A ₄	A ₅
В ₀	В ₁	В2	В ₃	В ₄	B ₅
c ₀	С ₁		С3	c ₄	C ₅
D ₀	D ₁	D_2	D ₃	D ₄	D ₅
E ₀	E ₁	E ₂	E ₃	E ₄	E ₅
F ₀	F ₁	F ₂	F ₃	F ₄	F ₅



A ₀	В ₀	c ₀	D ₀	E ₀	F ₀
A ₁	В ₁	с ₁	D ₁	E ₁	F ₁
A ₂	В2	c ₂	D ₂	E ₂	F ₂
A ₃	В ₃	С3	D ₃	E ₃	F ₃
A ₄	В ₄	c ₄	D ₄	E ₄	F ₄
A ₅	B ₅	C ₅	D ₅	E ₅	F ₅

Collective Operations are Blocking

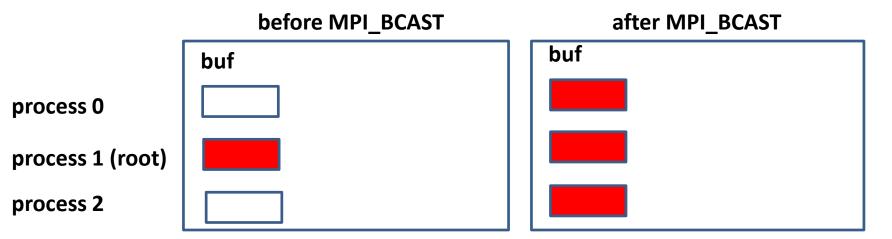


MPI_BARRIER: Synchronisation

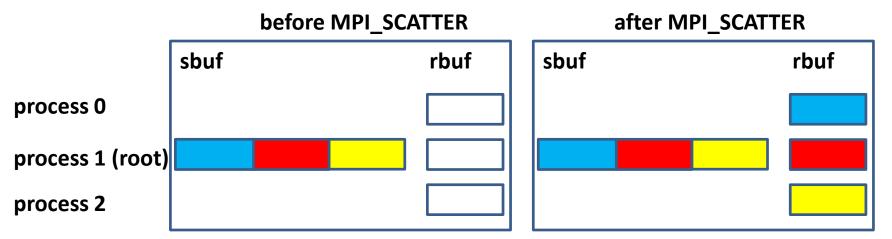
- MPI_BARRIER is usually not needed, because synchronization will be effected by other MPI routines
- MPI_BARRIER is useful for debugging and timing purposes

MPI_BCAST: Broadcast from root

- All processes must supply the same values for root and comm
- data_size = count * size (datatype)
 must be the same on all processes



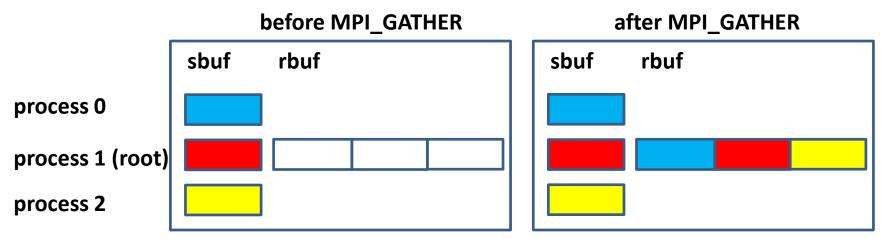
MPI_SCATTER: Scatter from root



Restrictions for Arguments in MPI_SCATTER

- All processes must supply the same values for root and comm
- r_data_size = rcount*size(rtype) on all processes
 must be equal to
 s_data_size = scount*size(stype) on process root
- **sbuf** is ignored an all non-**root** processes
- The total size of data scattered from process root is
 nproc * s_data_size

MPI GATHER: Gather to root



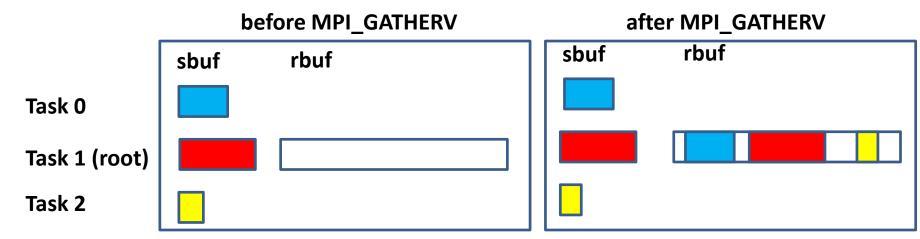
Restrictions for Arguments in MPI_GATHER

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 must be equal to
 r_data_size = rcount*size(rtype) on process root
- rbuf is ignored an all non-root processes
- The total size of data gathered on process root is
 nproc * r_data_size

Other Collective Communication Routines

- MPI_ALLGATHER similar to MPI_GATHER,
 - but all processes receive the result vector
 - > therefore no root argument
- MPI_ALLTOALL
 - > each process sends messages to all processes
- MPI_GATHERV, _SCATTERV, _ALLGATHERV, _ALLTOALLV
 - > collective communication routines with variable data layouts
 - > The counts of elements is different for each process,
 - different displacements of the element to be scattered from the send buffer resp. different displacements of the elements to be gathered in the receive buffer can be prescribed
 - ➤ Identical array of counts and array of displacements must be given as arguments in the call on all processes

MPI GATHERV: Gather to root

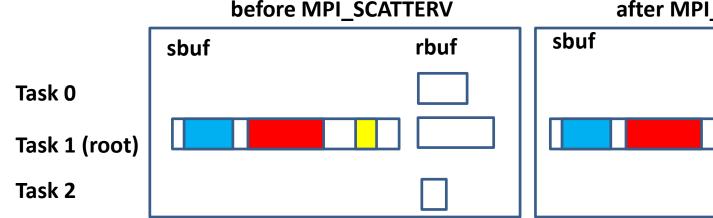


Restrictions for Arguments in MPI_GATHERV

The number of bytes in sbuf send from task i, determined by scount * size(stype),
must be equal to the number of bytes received in the i-th block in rbuf on root, determined by
rcounts(i) * size(rtype)

- The data block sbuf from task i will be stored in rbuf on root
 with displacement of displs(i) elements of type rtype from address rbuf.
- rbuf, rcounts, displs will be ignored on all non-root tasks.

MPI SCATTERV: Scatter from root





Restrictions for Arguments in MPI_SCATTERV

- The number of bytes in rbuf received on task i, determined by
 rcount * size(rtype),
 must be equal to the number of bytes sent from the i-th block in sbuf on
 root, determined by
 scounts(i) * size(stype)
- The i-th data block in sbuf has scounts(i) elements of type stype,
 is located at a distance of scounts(i) elements from the start address of
 sbuf and
 will be stored in the receive buffer rbuf on process i as rcount elements of
 type rtype
- sbuf, scounts, displs will be ignored on all non-root tasks.

In Place Variants

- In place variant of MPI_GATHER
 - The value MPI_IN_PLACE can be provided as argument for **sbuf** in the **root process**, if the root data to be gathered are already on their place in **rbuf** of the root
- In place variant of MPI_ALLGATHER
 - ➤ The value MPI_IN_PLACE can be provided as argument for **sbuf** in all processes, if the data to be gathered from the processes are already on their place in the **rbuf** of this process.

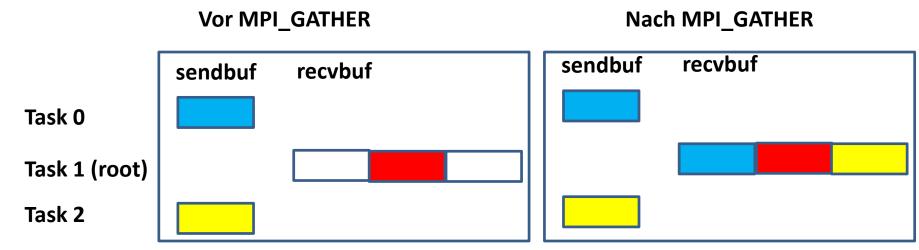
MPI_GATHER: mit MPI_IN_PLACE auf root

non-root-Task

MPI_GATHER(sendbuf, sendcount, sendtype, recvbuf, recvcount, recvtype, root, comm)

root-Task

MPI_GATHER(MPI_IN_PLACE, sendcount, sendtype, recvbuf, recvcount, recvtype, root, comm)



MPI ALLGATHER mit MPI IN PLACE

April 2024

MPI_ALLGATHER(MPI_IN_PLACE, sendcount, sendtype, recvbuf, recvcount, recvtype, comm)

Vor MPI_ALLGATHER Sendbuf recvbuf Task 0 Task 1 Task 2

Parallel Programming with MPI

26

Global Reduction

An example: Add the results local res computed in each of 3 tasks to a total result: total res = local res 0 + local res 1 + local res 2 INTEGER loc res, total res, all res(3) ! Gather local results into the array all res on the root process with MPI GATHER MPI GATHER (loc res, 1, MPI INTEGER , all res, 1, MPI INTEGER , root , comm, ierror) ! Add the elements of all res on the root process if (myid.eq.root) then total res = all res(1)+all res(2)+all res(3) end if

Global Reduction with MPI_REDUCE

Global Reduction Operations

To perform a global reduce operation across data on all processes of a communicator:

```
redo = d0 o d1 o d2 o d3 o ...

- di = data in process rank i

• single variable, or

• vector, i.e. di = (di1,di2,di3,...)

- o = associative operation

If di are vectors, the result of the reduce operation also is a vector: redo = ( d01 o d11 o d21 o d31 o ... ,

d02 o d12 o d22 o d32 o ... ,

d03 o d13 o d23 o d33 o ... ,

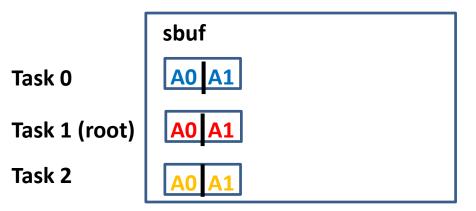
... )
```

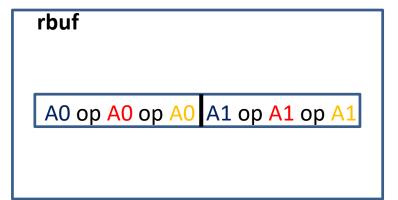
- Examples:
- global sum or product
- global maximum or minimum
- global user-defined operation

Predefined Reduction Operations

Meaning Name (fortran,c) (mpi4py) MPI MAX MPI.MAX maximum MPI MIN MPI.MIN minimum MPI SUM MPI.SUM sum MPI PROD MPI.PROD product MPI LAND logical and MPI.LAND MPI BAND bit-wise and MPI.BAND MPI LOR MPI.LOR logical or MPI BOR MPI.BOR bit-wise or MPI LXOR logical exclusive or (xor) MPI.LXOR MPI BXOR bit-wise exclusive or (xor) MPI.BXOR MPI MAXLOC MPI.MAXLOC max value and location MPI MINLOC MPI.MINLOC min value and location

 All processes must supply the same values for count, root, comm and datatype





Variants of Reduction Operations

- MPI_ALLREDUCE
 - returns the result in all processes
 - > no root argument
- MPI_REDUCE_SCATTER_BLOCK and MPI_REDUCE_SCATTER
 - result vector of the reduction operation is scattered to the processes into the result buffers
- MPI_SCAN
 - result at process with rank i:
 - =reduction of sbuf-values from rank 0 to rank i
- MPI_EXSCAN
 - result at process with rank i :
 - =reduction of sbuf-values from rank 0 to rank i-1

MPI_ALLREDUCE(sendbuf, recvbuf, count, datatype, op, comm)

Task 0

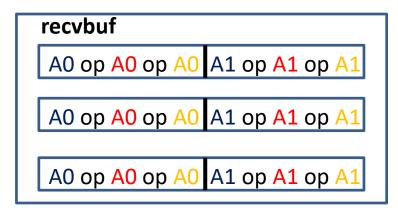
Task 1

Task 2

A0 A1

A0 A1

A0 A1



MPI_REDUCE_SCATTER(sendbuf, recvbuf, recvcounts, datatype, op, comm)

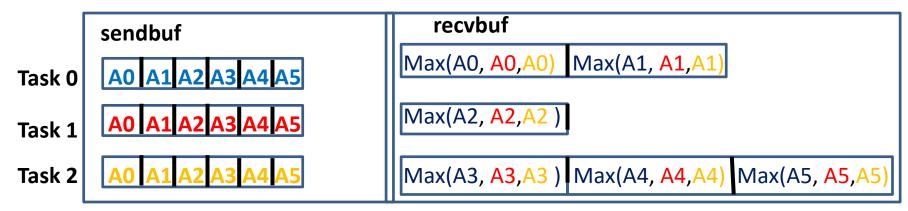
IN	sendbuf	starting address of send buffer (choice)
OUT	recvbuf	starting address of receive buffer (choice)
IN	recvcounts	non-negative integer array (of length group size)
		specifying the number of elements of the result
		distributed to each process.
IN	datatype	data type of elements of send and receive buffers
		(handle)
IN	ор	operation (handle)
IN	comm	communicator (handle)

MPI_REDUCE_SCATTER(sendbuf, recvbuf, recvcounts, datatype, op, comm)

Example: recvcounts(0) = 2, recvcounts(1) = 1, recvcounts(2) = 3 op = MPI_MAX

The number of elements in sendbuf to be reduced over nproc tasks is

recvcount(0)+ . . . +recvcount(nproc-1)



MPI_REDUCE_SCATTER_BLOCK(sendbuf, recvbuf, recvcount, datatype, op, comm)
z.B. recvcount = 2, op = MPI_MAX

The number of elements in sendbuf to be reduced over nproc tasks is

nproc*recvcount

Task 0

A0 A1 A2 A3 A4 A5

Task 1

Task 2

A0 A1 A2 A3 A4 A5

recvbufbuf

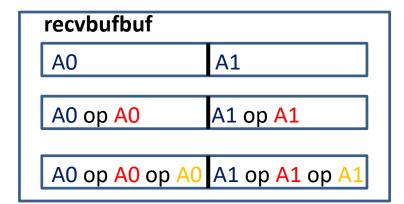
Max(A0, A0, A0) Max(A1, A1, A1)

Max(A2, A2, A2) Max(A3, A3, A3)

Max(A4, A4, A4) Max(A5, A5, A5)

MPI_SCAN(sendbuf, recvbuf, count, datatype, op, comm)

sendbuf
Task 0
A0 A1
Task 1
A0 A1
Task 2
A0 A1



MPI_EXSCAN(sendbuf, recvbuf, count, datatype, op, comm)

Task 0

A0 A1

Task 1

A0 A1

Task 2

A0 A1

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